Chapter 18
Demonic Fuzzy Relational Calculus

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ABSTRACT

In this chapter, the authors categorize methods that are used to formally specify and verify software requirements. They discuss several formal method-related subjects such as calculus fuzzy and relational calculus.

MOTIVATION

The importance of relations is almost self-evident. Science is, in a sense, the discovery of relations between observables. Zadeh has shown the study of relations to be equivalent to the general study of systems (a system is a relation between an input space and an output space.) (Goguen, 2009).

The calculus of relations has been an important component of the development of logic and algebra since the middle of the nineteenth century. George Boole, in his “Mathematical Analysis of Logic” (Boole, 1847), initiated the treatment of logic as part of mathematics, specifically as part of algebra. Quite the opposite conviction was put forward early this century by Bertrand Russell and Alfred North Whitehead in their Principia Mathematica (Whitehead & Russell, 1910): that mathematics was essentially grounded in logic. The logic is developed in two streams. On the one hand algebraic logic, in which the calculus of relations played a particularly prominent part, was taken up from Boole by Charles Sanders Peirce, who wished to do for the “calculus of relatives” what Boole had done for the calculus of sets, Peirce’s work was in turn taken up by Schröder in “Algebra und Logik der Relative.” Schröder’s work, however, lay dormant for more than 40 years, until revived by Alfred Tarski in his seminal paper “On the Calculus of Binary Relations” (Tarski, 1941). Tarski’s paper is still often referred to as the best introduction to the calculus of relations. It gave rise to a whole field of study, that of relation algebras, and more generally Boolean algebras with operators. This important work defined
much of the subsequent development of logic in the 20th century, completely eclipsing for some time the development of algebraic logic. In this stream of development, relational calculus and relational methods appear with the development of universal algebra in the 1930’s, and again with model theory from the 1950s onwards. In so far as these disciplines in turn overlapped with the development of category theory, relational methods sometimes appear in this context as well. It is fair to say that the role of the calculus of relations in the interaction between algebra and logic is now well understood and appreciated, and that relational methods are part of the toolbox of the mathematician and the logician.

The main advantages of the relational formalization are uniformity and modularity. Actually, once problems in these fields are formalized in terms of relational calculus, these problems can be considered by using formulae of relations, that is, we need only calculus of relations in order to solve the problems. This makes investigation on these fields easier.

Over the past twenty years relational methods have however also become of fundamental importance in computer science. For example, much of the theory of nonclassical logics is used (though sometimes re-invented) in the new so-called program logics. These arose from the realization that a program may be thought as an input-output relation over some state space: an accessibility relation. This point of view, that the calculus of relations is fundamental to programs, was clearly enunciated by the Oxford Group of Tony Hoare in an influential paper on “Laws of Programming” (Hoare, et al., 1987). Also during the 1980s, much of the equational theory of relation algebras were already being applied to program semantics and program development. For example, the book Relations and Graphs by Schmidt and Ströhlein (in German, 1989 [Schmidt & Ströhlein, 1989]; in English, 1993 [Schmidt & Ströhlein, 1993]) started from the basis of programs as graphs with Boolean matrices. Thus computer science, as the new application field for relational methods, has both drawn from and contributed to previous logic/mathematical work; this is a sign of healthy development.

**Demonic Relational Calculus**

In the context of software development, one important approach is that of developing programs from specifications by stepwise refinement (e.g. Back, 1981; Fchier, 2002a, 2003; Wirth, 1971). One point of view is that a specification is a relation constraining the input-output (respectively, argument-result) behaviour of programs. Inspired by interpretation of relations as non-deterministic programs, demonic, angelic and erratic variants of relational operations have been studied. The demonic interpretation of non-deterministic turns out to closely correspond to the concept of under-specification, and therefore the demonic operations are used in most refinement calculus (Wolfram, 2001).

The demonic calculus of relations (Boudriga, Elloumi, & Mili, 1992; Desharnais, Mili, & Nguyen, 1997) views any relation \( R \) from a set \( A \) to another set \( B \) as specifying those programs that terminate for all \( a \in A \) wherever \( R \) associates any values from \( B \) with \( a \), and then the program may only return values \( b \) for which \((a,b)\in R\). Consequently, a relation \( R \) refines another relation \( S \) if \( R \) specifies a larger domain of termination and fewer possibilities for return values. The demonic calculus of relations has the advantage that the demonic operations are defined on top of the conventional relation algebraic operations, and can easily and usefully be mixed with the latter, allowing the application of numerous algebraic properties.

Boudriga et al. (1992) give a refinement order introduced initially in Mili (1987). The set of relations with this order is a semilattice. Similar notions have been defined in Back (Back & von Wright, 1992; Back, 1981), Morgan and Robinson (1987), and Morris (1987). In the following,