

Efficient Identity-Based Multi-Cloud Security Access Control in Distributed Environments

S. K. Yakoob, Koneru Lakshmaiah Education Foundation, India*

V. Krishna Reddy, Koneru Lakshmaiah Education Foundation, India

ABSTRACT

Distributed computing is a forceful idea in disseminated registering which depicts versatile information to the executives for a minimal price dependent on client interest to various business associations. Because of multi-cloud identity-based encryption over distributed environment, in this document, the authors present and implement a novel identity-based multi-cloud security access control approach (NIMSACPA) for efficient security in multi data security and privacy based on three basic parametric concepts: 1) open minded security between autonomous user privacy using Byzantine protocol, 2) to classify the security privileges with respect to multi-cloud data sharing is described using DepSky Architecture, and 3) for identity-based information distribution between diverse users in CC described using Shamir secret key sharing procedure. This execution gives better and critical execution as far as data stockpiling and information investigation contrast and existing cryptographic techniques alongside practical multi-cloud information.

KEYWORDS

Access Control Policy, Byzantine Protocol, DepSky, Distributed Computing, Multi-Cloud Security, Shamir Secret Key Sharing

1. INTRODUCTION

Distributed computing is an arising and mature innovation in various business associations to diminish their association framework and adjust their association's answers for design cloud totally and to some degree. Various government and confidential associations keep up with complex foundations, for example, medical care/media communications are adjusting conveyed registering to lessen costs. Distributed computing has a few limits connected with security or protection in information rethinking between various clients in the cloud, which ought to contain and represent different complex ongoing applications. Cloud specialist co-op address protection/security issues as given the pressing/high need in information dissemination, security managing "Single Cloud" has become less force with clients in view of potential issues like help demand handling and accessibility regarding malignant clients present inside the single cloud. Secure distributed storage architecture supports the ability to transfer specific data using a limited intermediate re-encryption scheme. The encryption scheme

DOI: 10.4018/IJeC.316771

*Corresponding Author

This article published as an Open Access article distributed under the terms of the Creative Commons Attribution License (<http://creativecommons.org/licenses/by/4.0/>) which permits unrestricted use, distribution, and production in any medium, provided the author of the original work and original publication source are properly credited.

supports transmitting operations over encoded and encoded communications as well as decentralized elimination codes over scrambled messages. So as of late, moves towards security from a single cloud to “different mists” and “haze of mists.”

This chapter basically centers on security connect with conveyed climate with haze of mists. Information can be imparted to outsider cloud then, at that point, distributed computing clients dodges untrusted specialist co-op of the cloud to safeguard private or significant information like individual/MasterCard/subtleties from inside malevolent aggressors is a basic significance issue. Also, the potential for information security from single to haze of mists is analyzed in distributed computing. Scholar proposes a hybrid way to deal with empower security with respect to multi-cloud for proficient security in multi-mists, the half-breed method comprises three principal modules i. DepSky further develops respectability and secrecy of information put away in the cloud through encryption and unscrambling information on merit cloud from haze of entomb mists. ii. Byzantine convention to open-minded security breaks to server disappointments clouds which are free. iii. Shamir’s private sharing method further develops dependability and security of information stockpiling without decline/impact execution.

The fundamental commitments of this work are as follows:

1. This work tracks Byzantine adaptation to non-critical failure issues in multi-distributed computing to propose an adaptation to the internal failure model to portray proficiency of the haze of cloud climate.
2. Defines DepSky framework for secure and reliable information stockpiling framework, which depicts and fabricates business multi-cloud climate to haze of mists, and portray got to clients to summon individual client procedure on mists.
3. Present information on the executive’s model for the multi-cloud, which is Byzantine Fault Tolerance Multi-Cloud Database (BFTMCD) move which is a blend of Quantum Byzantine Agreement convention and Shamir’s Private Sharing way to deal with secure business information capacity in multi-cloud climate. Scholars contrast the proposed approach and existing cryptography models utilizing bury cloud climate.

2. REVIEW OF LITERATURE

Security angles connect with distributed computing with various creators/specialists. Likewise, talk about various dangers present in security of disseminated registering. As was previously discussed, cloud expert centers could assist clients, but security risks demand a crucial effort in the case of dispersed figuring (United Nations, 2012). When it comes to sharing of cloud data or framework workspaces, customers are misguided with risk of a loss of confidence (Syamsuddin, 2011). An ongoing IDC audit (Armbrust & Fox, 2010; Kim, 2009; Loshin, 2003; Nikkhah et al., 2011) shows that concerns about data security pose the greatest threat to the success of 74% of IT execution on CIOs of distributed figure appointments. Critical data must be shielded from outsiders and potentially malicious insiders, such as MasterCard subtleties or patients’ helpful records (Gâdescu, 2012). Moving data sets to a large server farm involve several security challenges, including issues with virtualization, accessibility, insurance and control of data obtained from intruders, dependability, protection, and data loss or theft. Data accumulation security, application security, data transfer security, similar fundamental concerns, and security about pariah resources are presented by Kim,W & Gâdescu,V et.al (2012).

The level of security commitment between customers and providers across various cloud benefit models is impressive. According to Amazon, customers are tasked with keeping an eye on the security of the company’s information technology infrastructure, which includes the servers, networks, workstations, and mobile devices that house the company’s operational systems, applications, and data, who claims that their EC2 watched out for security management in conjunction with physical, regular, and virtualized security. IaaS safety concerns include unique features; safety concerns have a more significant impact on

everyone's cloud than they do on private clouds. Cloud companies are online-based, so any problem with web security will also impact them. Cloud service provider bases security on the distributed environment, resources in the cloud is accessed over the web. As a result, data is now conveyed to the clients through web orchestration, which may be shaky. The cloud will also be affected by web security vulnerabilities. Additionally, because there are substantial resources reserved inside clouds and clouds are weak, cloud hazards are theoretically dangerous. The cloud's usage of technology is similar to that of the Internet. Insufficient encryption methods and secure networks are available to support data transmission in the cloud. Cloud conditions must be safe and private for users to prevent data interference from software developers and sophisticated criminals using the Internet.

NIST defines conveyed figuring as "a framework for enabling valuable, on-demand organize explore to such a prevalent & configurable handling resources (servers) which might be immediately reconfigured and published with irrelevant collaboration of experts present in organization". It is crucial to move widely dispersed numbers from a single fog to multiple fogs to ensure the safety of client data in transmission. The expression "multi-fogs" is similar to Vukolic's "between cloud" and "dimness of-fogs" articulations. Additionally, they advocate against performing distributed computation using a single cloud. A cloudy sky connects various fog types and tones by utilizing their structure, resulting in clear use and distinct areas. The multi-cloud condition, which manages a few fogs and avoids dependence on any one particular cloud, has received the most attention in continuous assessment. For specific circumstances, switching from internal or single clouds to many fogs makes sense. At this point, organizations of a single cloud are reliant upon a power interruption, as N. Santos et al. (2016) have shown. Vukolic (2010) hypothesizes that moving to an entomb cloud will be motivated by a desire to improve upon the single cloud's track record of reliability, trustworthiness, and security by using many cloud service providers to distribute these traits more widely. Also, strong appropriated storing, which employs a subset of Byzantine adaptation to non-critical failure (BFT) methodologies, has been prescribed by S. Bouchenak et.al (2013) to be used in multi-fogs. Different continuous assessments around there have fabricated shows between clouds.

3. BASIC PRELIMINARIES

This part portrays the relocation of dispersed registering concerning security in single to multi-cloud to guarantee the security of various clients in distributed computing. Fundamental examination present in this work center around multi-distributed computing climate, which controls and partitions and dodges based on cloud climate. Fundamental depictions of a multi-cloud environment with security are as follows.

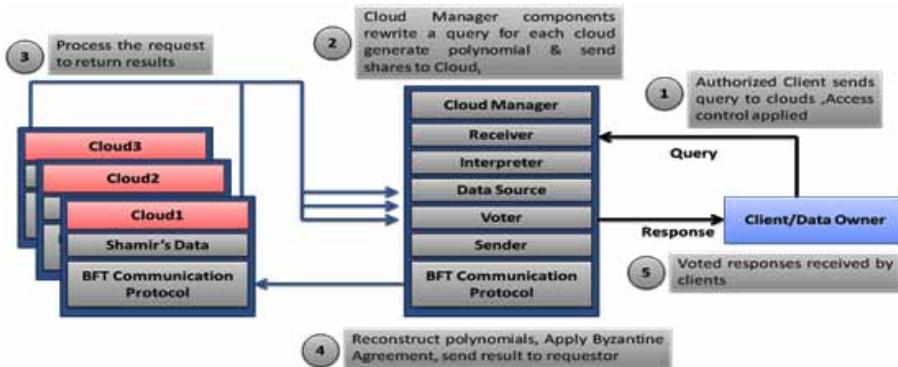
3.1 Byzantine Based Multi-Cloud Information the Executive Model

To plan a multi-cloud climate, for example, entomb cloud, Byzantine adaptation to non-critical failure is its concern. To sum up, Byzantine adaptation to non-critical failure involving Byzantine convention (Mohammed & AlZain, 2013) ordered progression in a basic and regular way. This model worked in light of the Quantum Byzantine Agreement convention and Shamir's Private Sharing way to deal with depicts Availability & Confidentiality, and honesty for continuous information sharing climate in multi-cloud. Byzantine Quantum Agreement Protocol (BQAP) Assuring reliable distributed computing that presents precipitation data transmitted over long distances is a bold notion. For this convention, there are two main cases:

1. All faultless processors have been combined into one.
2. The processor is faultless, giving it the same value regardless of its origin.

Figure 1 portrays information which created cloud data engineering, it characterizes multi-cloud board frameworks, and it depicts two circumstances:

Figure 1. Cloud information: The board model strategy



1. All non-flawed mists join with the same information.
2. If the cloud source is non-defective, in the event that information joins with the same information can be removed from unique cloud information.

3.2 Shamir Private Sharing Approach

It's the first defense against potential threats in a scenario with several cloud providers. Shamir's basic private-sharing (Goubin & Martinelli, 2011) method is illustrated as shown below.

Algorithm: Shamir Algorithm for generate multiple keys.

I/P: Sending documents as information.

O/P: Multi Secret Keys Ages for various enlisted clients for single document

Stage 1: Attaching cushioning piece of data, partition message into 64 pieces with products of 512 pieces.

Stage 2: Add the length (In double arrangement demonstrating length of the first message into 64 bit).

Stage 3: Start cushions sizes with comparable constants relying upon the quantity of words:

P=H0 = 0x67452301 Q=H1 = 0xEFCDAB89 R=H2= 0x98BADCFE

S=H3 = 0x10325476 T=H4 = 0xC3D2E1F0

Here each block size is viewed as 32 bit size for 160 pieces out put

Stage 4: Plan handling capabilities like

$$f(t;Q,R,S) = (Q \text{ AND } R) \text{ OR } ((\text{NOT } Q) \text{ AND } S) \quad (0 \leq t \leq 19)$$

$$f(t;Q,R,S) = Q \text{ XOR } R \text{ XOR } S \quad (20 \leq t \leq 39)$$

$$f(t;Q,R,S) = (Q \text{ AND } R) \text{ OR } (Q \text{ AND } S) \text{ OR } (R \text{ AND } S) \quad (40 \leq t \leq 59)$$

$$f(t;Q,R,S) = Q \text{ XOR } R \text{ XOR } S \quad (60 \leq t \leq 79)$$

Where "f" is non straight capability and "t" demonstrates handling round iterator.

Stage 5: Handling Message in 512 bit blocks: K(0), K(1), ..., K(79): 80 Handling Consistent Words, H0, H1, H2, H3, H4, H5: 5 Word cushions with starting qualities.

The algorithm portrays client information D into a number of capabilities (pieces) which the information k as for a number of pieces D_i which characterizes all-out information D . Complete information on $(k-1)$ pieces which isn't applicable to data about all-out information D and k ought to be not exactly consists n be the number of tokens, it keeps the upsides of shared keys that do not access information pieces. In view of this, Shamir's technique hypothetically secures data. The edge esteem (k,n) with $n=2k-1$ portrays pertinent and solid key administration situations to be accomplished. The principal objective of this portrays secure information base help as interaction to investigate private sharing strategy and offer various support suppliers, primary benefit of this approach characterizes the protection and security of reevaluated information.

3.3 DepSky System Architecture

This segment depicts about DEPSKY (Alysson, 2013) framework, it characterizes framework engineering, and afterward characterizes information and framework models with two created computations. Engineering for multi-cloud arrangement climate portrayed in Figure 2.

As shown in Figure 2, all the clouds in information sharing are stockpiling cloud to execute client code without limit, so clients get to their data with next to no progressions in the interface. DepSky calculations are carried out as a programming on client side in cloud. These library frameworks store UI comparative capacity in equal record handling frameworks, which permit read and write in backend information capacity. The heterogeneity of each cloud's specialized organizations needs to be explored using the DepSky library, and the findings of this research should be satisfied concerning the cloud in question. The information model shows the specifics of the new computations.

Figure 3 depicts the three distinct reflection levels that comprise the information model. In the first (left) box, we have a hypothetical data unit that is different from the primary storage item and on which the calculations are based (a register in conveyed figuring discourse). The information stored on a data unit object has a different name (X in the diagram), a transformation number (to aid in covering the article), confirmation information (in case of dispute, a secure hash based data representation), and a private key. In the ensuing focus, the applied data unit is carried out with unit in an imaginary cloud for storing purposes. There are two types of reports within a typical data compartment: metadata records for describing the data and the actual data records. The metadata files store the check data, the version

Figure 2. Fundamental design connects with DepSky for multi-cloud climate

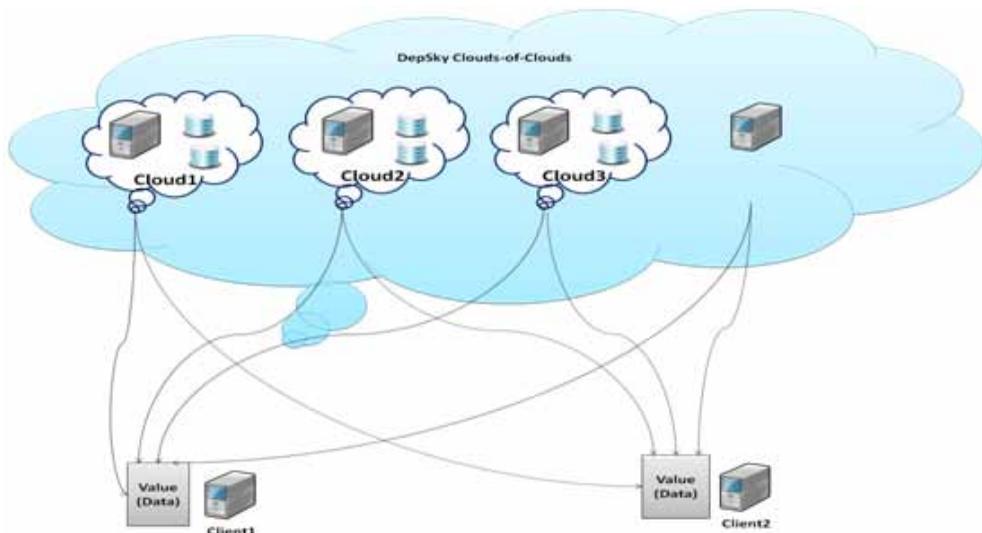
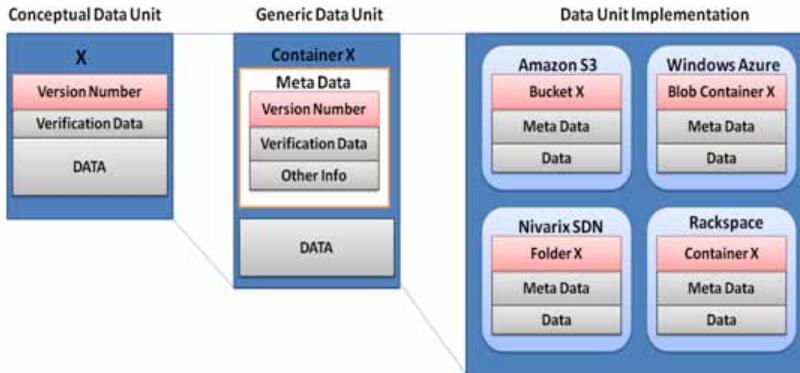


Figure 3. Data deliberation levels with DepSky



number, and other data that applications may need. Maintain similar data unit can store two heterogeneous data records, each with slightly different information. Data records are referred to as version of value, where version is the data's version number, and the metadata archive is called metadata (e.g., v1, v2, etc.). Thirdly, there is the execution of data units, which involves the holder being converted into the unique enhancements held by each cloud service (Bucket, Folder, etc.).

Multi-Cloud System Model depicts dispersed frameworks composed of three sorts of tasks, i.e., per users, journalists, and cloud storage specialist co-op displayed in Figure 3, while per users authors client jobs which are excessive with various cycles. Readers and scholars can fizzle erratic; all authors' per users of information unit d_u could share the key utilized to share a portion of the information composed on information unit while per users described in d_u can approach associated with key to confirm key marks.

Cloud server supplier is displayed with detached touchy information stockpiling with five tasks a) LIST (number of users b) explore (reads document in the distributed environment) c) make (make a holder) set (writers or altered the document). Cloud administration expects to give hierarchy in terms of access control to a framework to explore per users conjuring the rundown and get activities from capacity information in the cloud.

Algorithm: DepSky_AC implementation procedure

1. Pro (MultiSky_Write (du, Value)
2. Start
3. If $m_{V_{du}} = 0$
4. $N \leftarrow \text{query}_d(\text{du})$
5. $m_{V_{du}} \leftarrow m(\{n[i].v: 0 \leq i \leq l-1\})$
6. $\text{novel}_v \leftarrow m_{V_{du}} + 1$
7. $k \leftarrow \text{genSecKey}()$
8. $e \leftarrow E(v, k)$
9. $s[0 \dots l-1] \leftarrow s(k, l, f+1)$
10. $v[0 \dots l-1] \leftarrow s(e, l, f+1)$
11. for $(0 < i < l-1)$ do:
12. $d[a] = (s[a], e[b])$
13. $h[a] = H(d[a])$
14. $w[Q](\text{du}, 'v' + \text{new}_v, d)$
15. $n_m \leftarrow (n_v, h)$

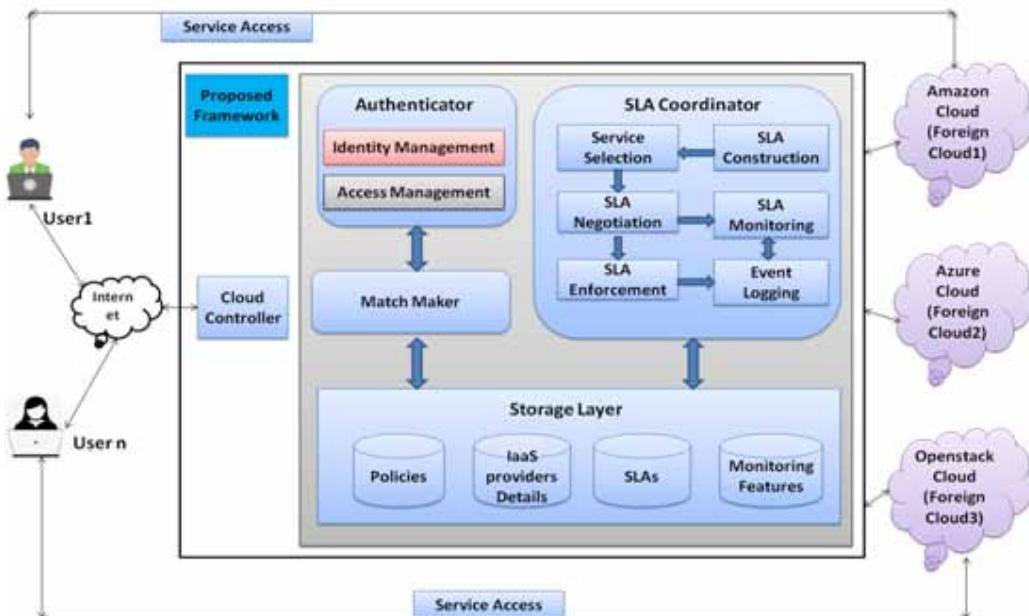
16. $s_i(n_m, K_i)$
17. $v[0..n-1] \rightarrow \text{new_meta}$
18. $w[Q](du, 'md', v)$
19. $m_{V_{du}} \rightarrow n_v$
20. $f(\text{DepSky_AC_Read}(du))$

4. PROPOSED APPROACH

When considering the less secure migration in cloud computing and the security notions described that, it is clear that a move from a single cloud to several clouds is necessary to guarantee the safety of user data. This research has centered on developing a multi-cloud environment that can regulate multiple clouds independently of one another. Scholar offers a Hybrid Approach, a synthesis of three core ideas for effectively managing inside attackers in a multi-cloud setting. This implemented method describes and focused on multi-cloud computing which controls and eliminates association of cloud. To solve efficient security concerns in a multi-cloud environment, proposed a Novel Identity-based Multi-Cloud Security Access Control Approach (NIMCSACA), which combines three main concepts to control inside attackers in a multi-cloud environment.

Using byzantine protocol sequences perform efficient operations with minimal conditions. DepSky describes architecture with respect to multi clouds based on different requirements. Proposed hybrid method explores access controllability between multiple users in multi cloud security. DepSky architecture describes cloud setup, protocol relates to byzantine data maintenance with different conditions which are authenticated by multi-cloud environment. Basic representation of multi- cloud architecture described in Figure 4. This figure describes server control cloud users communicate with accessibility of different users in distributed environment. Description of identify of different users via multi cloud accessibility is described in following sections.

Figure 4. Proposed Hybrid work for multi-cloud environment



Objective: Main objective of this hybrid method is to improve cloud client requests from cloud and evaluates services of users in cloud of cloud relations. Basic aim of multi cloud accessibility between different users ($a=0,1,\dots,m$) with respect to users requests ($b=1,2,\dots,N$) in their sequences of cloud as follows:

$$O(b, a) = M \sum_{j=1}^M \sum_{i=1}^N (R_{ij} - C_{ij})$$

Here M, N be the requests of cloud servers with different users and evaluate services possibilities in $R_{a,b}$ by users.

4.1 Initialization of Cloud Setup With Privacy Protocol

Protocol relates different service requests, attributes, setup parameters which are required for environment relates to multi clouds. Here server retrieve a service request in multi cloud setup, privacy protocol certifies user capabilities described by multi clouds. If protocol service is not available in cloud then privacy of user gives service of cloud based on third party sequences to check request of user performances. Below algorithm defines authentication protocol sequences with multiple attributes.

Algorithm: Procedure to select user requests in cloud

```

1.          BEGIN: Boot to enable multi-cloud collaboration
2.          while(the system is running)
3.          LC(Auth_Service)-> Check(Cert) //
Checking clients certificate
4.          if Valid (Cert):
5.              goto 17
6.          else:
7.              Auth_service->Send_request(Cert, ID)->TP //request
a new Certificate to TP
8.          end if
9.          for i=1 to n:
10.             LC->M_d(LC)->TPi
11.             TPi->Protocol(Cert)
12.             TPi->User_certificate(Cert)->LC
13.             LC(Auth_Service)->Receive(Cert)
14.          end for
15.          LC(Auth_service->Check(Cert)
16.          if Valid(Cert)
17.              wait (Request)
18.          else
19.              goto 7
20.          end if
21.          END
    
```

The algorithm shows that the authentication protocol hierarchy defines a multi-cloud environment with authentication in setup between defined client clouds. Dynamic communication between multi-clouds, each client cloud check authentication based on Single Sign-On authentication procedure by each local cloud server in Local Cloud (LC) authenticated with client cloud based on available resource

services. In this approach, we use Trusted Party (TP) for identification of identity provider which available on cloud data by user request holds a digital identity which is authenticated and authorized by granted service request in cloud proposed The proposed hybrid strategy is comprised of the following elements:

1. Offers a high degree of granularity in selecting services per request, taking into account the dynamic capabilities of many clouds.
2. Factors related to service quality used for precise matching based on available service request specifications from customers.

Response time, availability, dependability, memory, encryption/decryption time, and latency are only some of the quality of service needs that can be met by a hybrid approach, which also makes it suitable for a high volume of service requests.

Cloud servers are able to provide the needed services once they receive a request response from many customers.

4.2 Phase 1: User Based Quality of Service Request

The process of selecting a service begins on a cloud server, where the criteria for that service are entered. This user type may have higher memory requirements than the average cloud service request. Additionally, users pull the accessible server services into their client cloud. It's possible to see the request processing in several different ways, including the following.

The term “requirements” (RQ) refers to a group of standards for the quality of service:

$$R_Q = \{1; q_2; q_3; \dots q_n\}$$

S be the set of available services with different functionalities:

$$S = \{s_1; s_2; s_3; \dots s_n\}$$

Each service S has Q_s with different properties:

$$Q_s = \{Q_{s1}; Q_{s2}; Q_{s3}; \dots Q_{si}\}$$

where $Q_{si} = \{q_{i1}; q_{i2}; q_{i3} \dots q_{ij}\}$, which represents quality matrices of service requests.

4.3 Phase 2: Matrix Construction Based on Requirement

When a user has collected quality of service parameters, they can use them to compile a list of all the services that meet their criteria, along with the quality of service parameters for each one. This list is then used to construct an accuracy matrix that ranks the services according to how well they satisfy the user's requests:

$$R_Q = Q_{Si}$$

If there is no user Q_R in the quality service matrix Q_{Si} , then Q_{Si} is null and Q_R is 0. The number of requests R_Q that can be satisfied by the recognized services m is used to generate an accuracy matrix. To define the necessary matrix representation R, we can do the following:

$$\begin{matrix} S_1 \\ S_2 \\ \dots \\ S_m \end{matrix} \begin{pmatrix} R_{Q1} & R_{Q2} & \dots & R_{Qn} \\ r_{11} & r_{12} & \dots & r_{1n} \\ r_{21} & r_{22} & \dots & r_{2n} \\ \dots & \dots & \dots & \dots \\ r_{m1} & r_{m2} & \dots & r_{mn} \end{pmatrix}$$

If the mandatory requirements of the QoS are not fulfilled, they will be removed from the R_Q processing queue.

4.4 Phase 3: Accuracy of Matrix Representation

Accuracy of matrices calculated Ideal services, for instance, need service availability as well as throughput/memory metrics with response time to explore low latency, and this is true regardless of the QoS parameters used to measure latency. If the user goes with the high-latency, dependent requests below:

$$\frac{Q_{ij}}{Q_1} \text{ when } (Q_{ij} < Q_1)$$

$$\frac{Q_{ij} - Q_1}{Q_h - Q_1} + \alpha$$

$$\frac{Q_{ij} + \beta}{Q_{\max}} \text{ when } (Q_{ij} > Q_h)$$

For low latency values as follows:

$$\frac{Q_{ij}}{Q_1} \text{ when } (Q_{ij} < Q_h)$$

$$\frac{Q_{ij} - Q_1 + \alpha}{Q_h - Q_1} \text{ when } (Q_1 \leq Q_{ij} \leq Q_h)$$

$$\frac{Q_{\min}}{Q_{ij}} + \beta \text{ when } (Q_{ij} > Q_1)$$

Value of the i_{th} Quality of Service (QoS) feature of the j th service (represented by Q_{ij}), minimum range of user expectations (represented by Q_1), and maximum range of user expectations (represented by Q_h) are used in each of the aforementioned equations. Q_{\max} and Q_{\min} represent the highest and lowest possible scores on a Quality of Service (QoS) scale supplied by a service provider. And also fit in the interval [1, 2, 3,...], where $\alpha < \beta$, the results under the aforementioned criteria have been normalized to the interval [0, 1]. In a multi-cloud environment with many distinct sorts of cloud users, we advocate a solution based on the aforementioned process for authenticating users.

5. EXPERIMENTAL RESULTS

To calculate the efficiency of the implemented method design empirically in different clouds with multiple users. The experimental setup conducted in our implementation was performed to assess

the efficiency/scalability of the proposed approach with the decrease of runtime overheads with a collaborative multi-cloud environment. In the multi-cloud configuration that we've created, we deploy cloud environments built in JAVA with the NETBEANS tool and the CloudSim module. Each computer has a clock speed of 2.4 GHz, 8GB of RAM, and a storage capacity of 1TB. To circumvent strict security measures, both like user of cloud and service of cloud have turned to cloud web services. Our experimental configuration makes use of a cloud server with a high request throughput and several cloud client instances in each cloud client to ensure the scalability and efficiency of the suggested technique. This experiment aims to compare how well the suggested method performs against alternative methods of calculating the same set of operations in terms of how much time each takes. The following figures depict the amount of time spent by various instances of the client cloud during authentication using the suggested method. This research aims to assess the efficacy of the proposed hybrid strategy by comparing it to other authentication techniques, such as SAML (Memon, 2021), SAML with Proxy Re-encryption (Pachala et al., 2021), and Kerberos (n.d.).

Following figures show the presentation of implemented method by concerning set up time, time of encoding/encryption, unscrambling time, normal exactness in network precision and memory for client transferred demands as documents and client input subtleties in multi cloud information security. Table 1 addresses the ideal opportunity for various client cases information partaking in cloud.

As displayed in Table 1, contrasted with customary methodologies for example SAML, with intermediary set aside around equivalent effort to investigate client examples on cloud at whatever point client occasions expanded then those methodologies set aside more effort to execute the administrations of various clients. Where when contrasted with crossover approach, it was taken by less time contrast with existing methodologies.

Table 2 shows time of encryption with administration demands done by clients to transfer information in encoded design with various documentations via multi cloud storage environment.

Table 2, shows the encryption time assessment esteems and execution assessment of various methodologies in encryption, SAML, SAML with intermediary was taken additional time at whatever point increment the client case regarding various administrations. Proposed methodology was set aside less effort for encryption of documents transferred by various clients.

Decoding time with various client demand occasions to get information from multi distributed environment described in Table 3.

Table 3 shows the time assessment esteems and execution assessment of various methodologies in decoding likewise, SAML, SAML with intermediary was taken additional time at whatever point increment the client occasion concerning various administrations. Proposed approach was set aside less effort for decoding of documents transferred by various clients with various moment administrations.

Based on above figure show the exhibition as for complete time, encryption, decoding time and memory with various client examples, Proposed approach gives preferable execution over customary

Table 1. Time values based on user request

Different users	Time (Seconds)			
	SAML	SAML with Proxy	Kerberos	NIMCSACA
5	4.3	3.7	3.6	2.4
10	5.4	4.8	5.2	4.1
15	6.4	5.4	4.6	4.5
20	7.3	6.2	5.4	4.6
25	8.6	7.4	6.3	6.1
30	9	7.8	6.9	6.6

Table 2. Encryption time values with different user instances

Number of User Instances	Time (Seconds)			
	SAML	SAML with Proxy	Kerberos	NIMCSACA
10	4.3	3.7	3.6	3.5
20	5.4	4.8	5.2	4.1
30	6.4	6.7	5.2	4.3
40	7.3	6.2	7.1	4.6
50	8.6	7.4	6.3	7.3

Table 3. Decryption Time with different user instances

Number of User Instances	Time (Seconds)			
	SAML	SAML with Proxy	Kerberos	NIMCSACA
10	3.7	4.7	4.2	3.8
20	4.2	5.1	4.8	3.6
30	3.7	5.7	5.3	4.3
40	6.3	6.1	4.6	4.7
50	5.7	6.4	6.8	5.3

methodologies like SAML, SAML with Proxy Re-encryption and Kerberos planned with conditions of multi cloud environment.

6. CONCLUSION

Multi cloud offer a useful and promising solution for different cloud related framework with client benefits yet adaption of their data sharing security surrenders gets hardships nonappearance of wellbeing for appropriated structure. This goal familiarizes a proposed executed structure with handle secure elaboration over multi-cloud access organizations running consistently cloud specialist and client cloud. This hybrid methodology introduced reliant upon three particular circumstances i.e DepSky designing, the execution of client data is plagued by Byzantine. Shamir’s method for confirming data partitioning amongst many fogs. When compared to conventional methods of cloud data security, the proposed method achieves superior results in tests conducted on consistent cloud systems for cloud work. Test results performed on continuous cloud frameworks for cloud facilitating and consequences of the proposed approach give preferable outcomes over conventional methodologies utilized in cloud information security.

REFERENCES

- Alysson, B. M. C. (2013). DEPSKY: Dependable and Secure Storage in a Cloud-of-Clouds. *ACM Transactions on Storage*, 9(4).
- Armbrust, M., & Fox, A. (2010). A View of Cloud Computing. *Communications of the ACM*, 53(4), 50–58.
- Bouchenak, S., Chockler, G., Chockler, H., Gheorghe, G., Santos, N., & Shraer, A. (2013). Verifying cloud services: Present and future. *Operating Systems Review*, 47(2), 6–19.
- Gădescu, V. (2012). *Benefits of IPv6 in Cloud Computing* [Master of Science Thesis]. Tampere University of Technology.
- Goubin, L., & Martinelli, A. (2011). Protecting AES with Shamir's secret sharing scheme. *Springer International Workshop on Cryptographic*.
- Kerberos. (n.d.). <http://web.mit.edu/kerberos>
- Kim, W. (2009). Cloud Computing: Today and Tomorrow. *Journal of Object Technology*, 8(1).
- Loshin, P. (2003). *IPv6* (2nd ed.). The Morgan Kaufmann.
- Memon. (2021). *Federated Access to Collaborative Data and Compute Infrastructures* [PhD dissertation]. University of Iceland.
- Mohammed, A., & AlZain, B. S. (2013). A Byzantine Fault Tolerance Model for a Multi-Cloud Computing. *IEEE 16th International Conference on Computational Science and Engineering*.
- Nikkhah, M., Gu'erin, R., Lee, Y., & Woundy, R. (2011). Assessing IPv6 through web access a measurement study and its findings. *Proc. ACM CoNEXT*, 26:1–26:12.
- Pachala, S., Rupa, C., & Sumalatha, L. (2021). An improved security and privacy management system for data in multi-cloud environments using a hybrid approach. *Evolutionary Intelligence*, 14, 1117–1133. doi:10.1007/s12065-020-00555-w
- Santos, N., Pereira, S., Alves, A., & Chaves, R. (2016). Storekeeper: A security-enhanced cloud storage aggregation service. In *2016 IEEE 35th Symposium on Reliable Distributed Systems (SRDS)* (pp. 111-120). IEEE.
- Syamsuddin, I. (2011). Evaluation of E-government Initiatives in Developing Countries: An ITPOSMO Approach. *International Research Journal of Applied and Basic Sciences*, 2(12), 439–446.
- United Nations. (2012). *E-Government Survey 2012 E- Government for the People*. United Nations.
- Vukolic, M. (2010). The Byzantine empire in the intercloud. *ACM SIGACT News*, 41, 105–111.

Sk. Yakob received M.Tech degree in Computer Science and Engineering from Anurag Engineering College, Jawaharlal Nehru Technological University, Hyderabad, Telangana, India in 2009. He is presently working as Associate Professor in Dept. of CSE, Sai Spurthi Institute of Technology, Sathupally. He is pursuing PhD in KL University, Vijayawada. His area of interest includes Cloud Computing and Information Security.

V. Krishna Reddy is presently Vice Principal & Professor in the Department of Computer Science & Engineering, KL University-Vijayawada, Andhrapradesh, India. He received PhD degree from Acharya Nagarjuna University, Guntur, Andhrapradesh. His research interests include Cloud Computing, network security, web services and Data mining. Dr. V. Krishna Reddy published more than 20 papers in refereed international journals and 15 papers in conferences and has been involved many international conferences as Technical Chair and tutorial presenter. He is an active member of IEEE, ISTE and Computer Society India.