Database Replication and Clustering for High Availability

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INTRODUCTION

The cloud-computing trend today has pushed ever more services that span virtually all business and government sectors to hundreds of millions of users. Users typically take it for granted that such services would be continuously available, 24 hours a day, 7 days a week. Any extended disruption in services, including both planned and unplanned downtime, can result in significant financial loss and negative social impacts. Consequently, the systems providing these cloud services must be made highly available.

A cloud service is typically powered by a multi-tier system, consisting of Web servers, application servers, and database management systems, running in a server farm environment. The Web servers handle direct Web traffic and pass requests that need further processing to the application servers. The application servers process the requests according to the predefined business logic. The database management systems store and manage all mission critical data and application state so that the Web servers and application servers can be programmed as stateless servers. (Some application servers may cache information, or keep session state. However, the loss of such state may reduce performance temporarily or may be slightly annoying to the affected user, but not critical.) This design is driven by the demand for high scalability (to support large number of users) and high availability (to provide services all the time). If the number of users has increased, more Web servers and application servers can be added dynamically. If a Web server or an application server fails, the next request can be routed to another server for processing.

Inevitably, this design increases the burden and importance of the database management systems. However, this is not done without good reasons. Cloud applications often access and generate huge amount of data on requests from large number of users. A database management system can store and manage the data in a well-organized and structured way (often using the relational model). It also provides highly efficient concurrency control on accesses to shared data.

While it is relatively straightforward to ensure high availability for Web servers and application servers by simply running multiple copies in the stateless design, it is not so for a database management system, which in general has abundant state. The subject of highly available database systems has been studied for more than two decades and there exist many alternative solutions (Agrawal et al., 1997; Cecchet, Candea, & Ailamaki, 2008; Drake et al., 2005; García, Rodrigues, & Preguiça, 2011; Kemme & Alonso, 2000; Patino-Martínez et al., 2005). In this article, we provide an overview of two most popular database high availability strategies, namely database replication and database clustering. The emphasis is given to those that have been adopted and implemented by major database management systems (Banker, 2011; Davies & Fisk, 2006; Vallath, 2004).

BACKGROUND

A database management system consists of a set of data and a number of processes that manage the data. These processes are often collectively referred to as database servers. The core programming model used in database management systems is called transaction processing. In this programming model, a group of read and write operations on the same data set are demarcated within a transaction. A transaction has the following ACID properties (Gray & Reuter, 1993):

- Atomicity: All operations on the data set agree on the same outcome. Either all the operations succeed (the transaction commits), or none of them are (the transaction aborts).
- **Consistency**: If the database is consistent at the beginning of a transaction, then the database remains consistent after the transaction commits.
- **Isolation**: A transaction does not read or overwrite a data item that has been accessed by another concurrent transaction.
- **Durability**: The update to the data set becomes permanent once the transaction is committed.

An example of an atomic transaction is shown in Figure 1. This transaction involves a debit operation on a savings account and a credit operation on a checking account. If both operations are successful and the user decides to commit the transaction, the changes to both accounts are made permanent, as shown in Figure 1(a). On the other hand, if the user decides to abort the transaction, the changes to both accounts will be reversed so that the account balances of both accounts are restored to the values at the beginning of the transaction, as illustrated in Figure 1(b).

To support multiple concurrent users, a database management system uses sophisticated concurrency control algorithms to ensure the isolation of different transactions even if they access some shared data concurrently (Bernstein et al., 1987). The strongest isolation can be achieved by imposing a serializable order on all conflicting read and write operations of a set of transactions so that the transactions appear to be executed sequentially. Two operations are said to be conflicting if both operations access the same data item and at least one of them is a write operation, and they belong to different transactions. Another popular isolation model is the snapshot isolation. Under the snapshot isolation model, a transaction performs its operations against a snapshot of the database taken at the start of the transaction. The transaction will be committed if the write operations do not conflict with any other transaction that has committed since the snapshot was taken. The snapshot isolation model can provide better concurrent execution than the serializable isolation model.

A major challenge in database replication, the basic method to achieve high availability, is that it is not acceptable to reduce the concurrency levels. This is in sharp contrast to the replication requirement in some other field, which often assumes that the replicas are single-threaded and deterministic (Castro & Liskov, 2002).

To achieve high availability, a database system must try to maximize the time to operate correctly without a fault and minimize the time to recover from a fault. The transaction processing model used in database management systems has some degree of fault tolerance in that a fault normally cannot corrupt the integrity of the database. If a fault occurs, all ongoing transactions will be aborted on recovery. However, the recovery time would be too long to satisfy the high availability requirement. To effectively minimize the recovery time, redundant hardware and software must be used. Many types of hardware fault can in fact be masked. For example, power failures can be masked by using redundant power supplies, and local communication system failures can be masked by using redundant network interface cards, cables and switches.

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**Figure 1. An example of atomic transactions**

(a) Transaction Manager | Savings Account | Checking Account
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Begin Transaction | Initial Balance: $5000 | Initial Balance: $1000
Withdraw $1000 | | |
Balance: $4000 | Deposit $1000 | |
Balance: $2000 | Make changes permanent | |
Final Balance: $4000 | Final Balance: $2000 | |
Commit Transaction

(b) Transaction Manager | Savings Account | Checking Account
--- | --- | ---
Begin Transaction | Initial Balance: $5000 | Initial Balance: $1000
Withdraw $1000 | | |
Balance: $4000 | Deposit $1000 | |
Balance: $2000 | Rollback changes | |
Final Balance: $4000 | Final Balance: $2000 | |
Abort Transaction