Consistency Is Not Enough in Byzantine Fault Tolerance

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INTRODUCTION

In Byzantine fault tolerance (BFT), a core concern is to ensure the consistency of replicas despite malicious attacks from the clients and compromised replicas (Zhao, 2014). This is accomplished by totally ordering incoming requests and by rendering the replica’s operations deterministic (Zhang et al., 2011). In the presence of application nondeterminism, such as the access of local clocks, replicas are rendered deterministic by forcing all non-faulty replicas to use the same values either supplied by the primary or computed deterministically. While this approach works well for some applications, such as a replicated file system, doing so could lead to the compromise of the service integrity for applications that rely on the use of random numbers.

For example, consider an Internet application that relies on the use of session-ids for stateful interactions between the server and its clients. As pointed out in (Dorrendorf, Gutterman, & Pinkas, 2007), if the session-id of an active session can be predicted, the client’s session with the server could be hijacked, which could lead to the leak of confidential information regarding the client, such as name, address, and the order history, or unauthorized orders (if the one-click option for placing orders is enabled). The session-id may be predicted by searching the limited entropy space if weak random bits are used in an application. For example, the authors of (Dorrendorf, Gutterman, & Pinkas, 2007) reverse-engineered a version of Tomcat (a popular Java Servlet Engine) and the related operations in a Window’s based Java Virtual Machine. They could attack the system by performing about 251 searches in finding an active session-id.

Therefore, it is critical not to weaken the strength of the random bits essential for the integrity of their operations when replicating these systems for Byzantine fault tolerance. For a sound coordination algorithm, it is essential to enable each replica to access its own entropy source and maintain its independence in such operations. However, this desire is in conflict with the basic requirement for state machine replication (Schneider, 1990), which mandates that the replicas must be deterministic or rendered deterministic to maintain strong replica consistency. The conflicting requirements for security and replication must be reconciled to avoid the dilemma of either favoring security over high availability by not performing state machine replication of the systems, or trading security for high availability by removing the randomness of the systems in order to perform state machine replication.

In this chapter, we present a novel replica coordination algorithm, referred to as the Collective-Determination BFT algorithm, or CD-BFT algorithm in short, towards the reconciliation of the conflicting requirements for security and for strongly consistent replication. The central idea behind this algorithm is that all random numbers to be used by the replicas are collectively determined, and furthermore, the determination is based on the contributions made by a quorum of replicas, at least one of which is correct.

In the CD-BFT algorithm, the replicas first reach a Byzantine agreement on the set of contributions from replicas, and then apply a deterministic algorithm, such as the bitwise exclusive-or.
operation (Young & Yung, 2004), to compute the final random value. The freshness of the random numbers generated is application dependent. Our approach does not alter the freshness of the random numbers. If a pseudo-random number generator is used, it should be periodically reseeded from a good entropy source.

BACKGROUND

An arbitrary fault is often referred to as a Byzantine fault. The term was introduced in (Lamport, Shostak, & Pease, 1982) to highlight a specific faulty behavior that a Byzantine faulty process may disseminate conflicting information to other processes. For example, a compromised process might exhibit such Byzantine faulty behavior. Byzantine fault tolerance refers to the capability of tolerating Byzantine faults in a system. It can be achieved by introducing sufficient redundancy into the system and by using a sophisticated replica coordination algorithm that can cope with Byzantine faulty replicas and clients. A basic requirement for such an algorithm is to ensure that all server replicas agree on the total ordering of the requests received despite the existence of Byzantine faulty replicas and clients. Such an agreement is often referred to as Byzantine agreement (Lamport, Shostak, & Pease, 1982).

Recently, a number of efficient BFT algorithms (Castro & Liskov, 2002; Kotla et al., 2007; Yin et al., 2003) have been proposed. Our CD-BFT algorithm is derived from the PBFT algorithm and we use the same system model as that in (Castro & Liskov, 2002). The PBFT algorithm operates in an asynchronous distributed environment. The safety property of the algorithm, i.e., all correct replicas agree on the total ordering of requests, is ensured without any assumption of synchrony. However, for the algorithm to make progress towards a Byzantine agreement (i.e., liveness), certain synchrony is needed, for example, a reasonable assumption is that the message transmission and processing delay has an asymptotic upper bound. This bound is dynamically explored in the algorithm in that each time a view change occurs, the timeout for the new view is doubled.

How to ensure strong replica consistency in the presence of replica non-determinism has been of research interest for a long time (Castro & Liskov, 2002; Castro, Rodrigues, & Liskov, 2003; Powell, 1991; Sleember & Narasimhan, 2006), most of which are for fault tolerant systems using the benign fault model. However, while the importance of the use of good random numbers has long been recognized in building secure systems (Viega & McGraw, 2002), we have yet to see substantial research work on how to preserve the randomized operations necessary to ensure the system integrity in a fault tolerant system. For the type of systems where the use of random numbers is crucial to their service integrity, the benign fault model is obviously inadequate and the Byzantine fault model must be employed if fault tolerance is required.

Some form of replica non-determinism, in particular those related to timestamp operations, has been studied in the context Byzantine fault tolerant systems (Castro & Liskov, 2002; Castro, Rodrigues, & Liskov, 2003). However, as we will argue in the next section, the existing approach is vulnerable to the presence of colluding Byzantine faulty replicas and clients. We also studied the replica non-determinism issue with the emphasis of classification of non-determinism types and the systematic handling of various types of replica non-determinism (Zhang et al., 2011).

RECONCILIATION OF SECURITY AND REPLICATION REQUIREMENTS

Pitfalls in Controlling Replica Randomness

In this section, we analyze a few well-known approaches that one may attempt to use to ensure replica consistency in the presence of replica randomness. We show that they are not robust